

Lec 6. M126

guess n=2 WS.

① 1/24/12

HW2: for log plots, easiest to use semilog, loglog, etc rather than take log of data is ln not log<sub>10</sub>.

demo importance of Chebyshev vs equally-spaced nodes!

recall quadrature  $Q_n(f) = \sum_{j=0}^n w_j f(x_j)$   $\exists \{w_j\}$  st.  $Q_n$  exact  $\forall f \in P_n$ .

last time: Interpolation quadrature on  $[a, b]$ ,  $n=1$ , ie  $n+1=2$  nodes, choose  $x_0=a, x_1=b$ , get  $Q_1(f) = h \frac{f(a)+f(b)}{2}$  trapezoid rule.

Thm (2.1): Let  $f \in C^2[a, b]$ , then  $|\int_a^b f(x) dx - Q_1(f)| \leq \frac{1}{12} \|f''\|_{\infty} h^3$

Pf [Thm (2.1) LIE]: Peano kernel  $K(x) = \frac{1}{2}(x-a)(b-x) \geq 0$  on  $[a, b]$ ,  $K'' = -1$  on  $[a, b]$ .  
 $\int_a^b K(x) f''(x) dx = -\int_a^b K'(x) f'(x) dx + [K(x) f'(x)]_a^b = -\int_a^b K''(x) f(x) dx - [K'(x) f(x)]_a^b = -\int_a^b f(x) dx + \frac{h}{2}(f(a)+f(b))$   
 magnitude  $\leq \int_a^b K(x) dx \cdot \|f''\|_{\infty} \frac{h^2}{6}$  QED.

Many split interval into smaller & apply above to each:  $\frac{a}{h} \frac{h}{h} \dots \frac{b}{h}$  "composite trapezoid rule"

error  $\leq \frac{1}{12} \|f''\|_{\infty} h^3 \cdot \frac{b-a}{h} \cdot \frac{1}{h}$  # intervals  
 $= \frac{b-a}{12} \|f''\|_{\infty} h^2 = O(h^2)$  convergence algebraic, order = # of nodes  $(n+1)$

Can increase n: more points on single interval, eg  $n=2$  Simpsons' (1743)

Say choose  $n$  equispaced pts.  $\frac{a}{h} \dots \frac{b}{h}$  single interval. Guesses for  $w_j$  as  $n \rightarrow \infty$ ?  $w_j$  are  $\int_a^b \delta_j(x) dx \rightarrow$  exp. large & oscillatory  $\rightarrow$  bad for roundoff.

Another way in which negative weights bad: Convergence.

Consider seq.  $(Q_n)_{n=0}^{\infty}$  of schemes.  $Q_n(f) := \sum_{j=0}^n w_j^{(n)} f(x_j^{(n)})$   
 Defn  $(Q_n)$  conv. if  $Q_n(f) \rightarrow Q(f)$  as  $n \rightarrow \infty$ ,  $\forall f \in C[a, b]$  nice property! (Recall impossible for interp! Low error)  
 Thm (Szegő)  $(Q_n)$  conv.  $\iff (Q_n)$  conv. for all polys &  $\exists C$  st.  $\sum_{j=0}^n |w_j^{(n)}| \leq C$  th

note: these means if weights blow up as  $n \rightarrow \infty$ , cannot be conv! (egs. equispaced)

Facts 1)  $P$  = polys 'dense' in  $C[a, b]$ , meaning:  $\forall f \in C[a, b]$  &  $\forall \epsilon > 0$ ,  $\exists p \in P$  st.  $\|f-p\|_{\infty} \leq \epsilon$   
 2) each  $Q_n$  is lin. op:  $C[a, b] \rightarrow \mathbb{R}$  w/  $|Q_n(f)| \leq \|f\|_{\infty} \sum_{j=0}^n |w_j^{(n)}| \leq C \|f\|_{\infty}$   
 Weierstrass.

Pf: use facts 1 & 2 w/ Banach-Steinhaus thm; so in common, this is  $\|Q_n\|_{\infty}$ .  
 Let  $(Q_n)$  be seq. of bnd lin. ops, to  $\mathbb{Q}$  bnd lin. op,  $X = C[a, b]$  Banach space. operator norm.

pointwise convergence  $\iff (Q_n)$  uniformly bndd & convergent on dense subset  
 $\forall f \in X, \lim_{n \rightarrow \infty} \|Q_n f - Q f\| = 0$   $\iff \lim_{n \rightarrow \infty} \|Q_n - Q\| = 0$

Banach-Steinhaus is a variant of 'Principle of Uniform Boundedness'. both std. in ② 1/29/12  
 This is pretty abstract,  $\hookrightarrow$  the  $\{E\}$  in B.S. func. anal.

so let's prove the non-B.U.B. part, the  $\{E\}$  in Thm: (ptwise conv.  $\Leftarrow$  dense & unit ball)

For any  $\epsilon > 0$ ,

$$Q_n f - Q f = Q_n f - Q_n p + Q_n p - Q p + Q p - Q f$$

can find  $p \in P$  st.  $\|p - f\|_\infty \leq \epsilon$

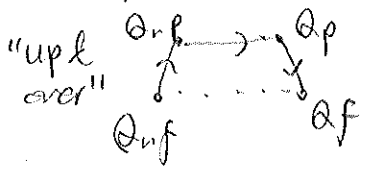
Take abs val & use tri. ineq:

$$|Q_n f - Q f| \leq \underbrace{|Q_n f - Q_n p|}_{\leq C \|f - p\|_\infty = C \epsilon} + \underbrace{|Q_n p - Q p|}_{\leq \epsilon} + \underbrace{|Q p - Q f|}_{\leq (b-a) \epsilon}$$

$\leq (C+1+b-a) \epsilon$  fixed const.

$\forall n > N$  for some  $N$  suff. large.  $\hookrightarrow$  if you like,  $\|Q\|$ .

So for each  $\delta > 0$ , choose  $\epsilon = \frac{\delta}{C+1+b-a}$  &  $\exists N$  st.  $\forall n > N$ ,  $|Q_n f - Q f| < \delta$ . QED.



bounded by  $\leq$  sum of 3 parts.

Also called " $\epsilon/3$ " argument, common in func. anal.

Positive weights is enough:

Corollary (9.11, Steklov): if  $(Q_n)$  conv for all polys,  $L^v w_j^{(n)} \geq 0$ , then  $(Q_n)$  conv.

pf:  $\|Q_n\|_\infty = \sum_{j=0}^n |w_j^{(n)}| = \sum_{j=0}^n w_j^{(n)} = Q_n(1) \xrightarrow[\text{poly}]{\text{fix } a} Q(1) = \int_a^b 1 dx = b-a$

nonneg.

so  $\exists C$  st.  $\|Q_n\|_\infty \leq C \forall n$ , use Szegö.

- Also minimal run-off error since sizes of weights as small as poss.
- Eg  $\Rightarrow$  composite trap. conv. (eve w's, conv.  $\forall$  polys since each has  $\|p''\|_\infty < \infty$ ). But equispaced

Now improved quadr. scheme on  $[a,b]$  ... w/ positive weights ...

Gaussian Quadrature (p9.3): optimal choice of nodes  $\rightarrow$  queen of quadratures on  $[a,b]$

$\rightarrow$  vrs. straight m.

but  $n=2$  degree 5 exact, generally can do degree  $2n+1$  exact, compared to only  $n$  for Newton-Lots of general nodes.  $\hookrightarrow$  defines Gaussian quadr w/ n nodes.

Let's show why:

• Orthogonality for fms.  $f \perp g \Leftrightarrow 0 = (f,g) := \int_a^b f(x)g(x) dx$

Lemma (9.13) Let  $x_0, \dots, x_n$  be distinct nodes of a Gaussian quadr.

Then  $q_{n+1}(x) := \prod_{j=0}^n (x-x_j) \perp p \quad \forall p \in P_n$  vanish at nodes.

pf:  $q_{n+1} p \in P_{2n+1}$  so  $\int_a^b q_{n+1}(x) p(x) dx \stackrel{\text{by Gauss.}}{=} \sum_{k=0}^n w_k q_{n+1}(x_k) p(x_k) = 0$  QED.